# On orthogonal generalized equitable rectangles

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#### Abstract

In this note, we give a complete solution of the existence of orthogonal generalized equitable rectangles, which was raised as an open problem in [4].

Key words: orthogonal latin squares, orthogonal equitable rectangles,

### 1 Introduction

A latin square of order t is a  $t \times t$  array defined on t symbols such that every symbol occurs exactly once in each row and exactly once in each column. Two latin squares of order t, say  $A = (a_{i,j})$  and  $B = (b_{i,j})$ , are orthogonal if the  $t^2$  pairs  $(a_{i,j}, b_{i,j}), 1 \le i \le t, 1 \le j \le t$ , are distinct.

Suppose  $r \leq t$ . An  $r \times t$  latin rectangle is an  $r \times t$  array defined on t symbols such that every symbol occurs exactly once in each row and at most once in each column. Two  $r \times t$  latin rectangles, say  $A = (a_{i,j})$  and  $B = (b_{i,j})$ , are orthogonal if the rt pairs  $(a_{i,j}, b_{i,j})$ ,  $1 \leq i \leq r$ ,  $1 \leq j \leq t$ , are distinct. It is easy to see that orthogonal  $t \times t$  rectangles are the same as orthogonal latin squares

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of order t. Orthogonal latin squares and orthogonal latin rectangles are well-studied combinatorial objects (see, e.g., [1]).

Stinson introduced orthogonal equitable rectangles in a recent paper [4]. Orthogonal equitable rectangles were motivated by a cryptographic application described in [3]. In fact, orthogonal equitable rectangles are a natural variation of orthogonal latin rectangles. An open question in [4] asked for necessary and sufficient conditions for the existence of a certain generalization of orthogonal equitable rectangles, which we define now.

Suppose  $r, t, s_1, s_2$  are positive integers such that  $rt = s_1s_2$ . Orthogonal generalized equitable rectangles (OGER) are defined to be two  $r \times t$  rectangles, say A and B, satisfying the following properties:

- 1.  $A = (a_{i,j})$  is defined on a set  $S_1$  of  $s_1$  symbols and  $B = (b_{i,j})$  is defined on a set  $S_2$  of  $s_2$  symbols, where  $s_1s_2 = rt$ .
- 2. A is equitable on rows and equitable on columns: each of the  $s_1$  symbols in  $S_1$  appears  $\lceil \frac{t}{s_1} \rceil$  or  $\lfloor \frac{t}{s_1} \rfloor$  times in every row of A, and  $\lceil \frac{r}{s_1} \rceil$  or  $\lfloor \frac{r}{s_1} \rfloor$  times in every column in A.
- 3. *B* is equitable on rows and equitable on columns: each of the  $s_2$  symbols in  $S_2$  appears  $\lceil \frac{t}{s_2} \rceil$  or  $\lfloor \frac{t}{s_2} \rfloor$  times in every row, and  $\lceil \frac{r}{s_2} \rceil$  or  $\lfloor \frac{r}{s_2} \rfloor$  times in every column in *B*.
- 4. A and B are orthogonal: the rt pairs  $(a_{i,j}, b_{i,j}), 1 \le i \le r, 1 \le j \le t$  are all distinct.

We denote A and B as  $(r, t; s_1, s_2)$ -OGER.

**Example 1.1** A (2,6;3,4)-OGER:

1	2	2	3	3	1	1 2	2	1	2	3	
2	3	3	1	1	3	3 4	1	2	1	4	

An  $(r, t; s_1, s_2)$ -OGER is a generalization of a pair of orthogonal equitable rectangles, which are discussed in [4]. In fact, an (r, t; r, t)-OGER is the same thing as a pair of orthogonal equitable  $r \times t$  rectangles. Furthermore, an (r, r; r, r)-OGER is just a pair of orthogonal latin squares of size r.

Stinson [4] gave an almost complete solution for the existence of orthogonal equitable rectangles. His solution only had a few possible exceptions, which were subsequently removed by Guo and Ge [2]. The following theorem summarizes these existence results.

**Theorem 1.2** There exists an (r, t; r, t)-OGER (i.e., a pair of orthogonal equitable  $r \times t$  rectangles) if and only if  $(r, t) \notin \{(2, 2), (2, 3), (3, 4), (6, 6)\}$ .

When  $\{r,t\} \neq \{s_1, s_2\}$ , orthogonal generalized equitable rectangles have no obvious cryptographic applications. However, their construction is a natural and interesting new problem in combinatorial designs. This problem at first glance seems difficult due to its generality:  $r, t, s_1, s_2$ can be any positive integers that satisfy the equation  $rt = s_1s_2$ . Despite the generality of the problem, we are able to completely solve it, using the result of Theorem 1.2 as a starting point, by applying three recursive constructions and three constructions of OGERs for individual parameter sets. The resulting solution is remarkably short.

#### 2 Main Theorem

In this section, we prove our main theorem. We begin by stating two lemmas that indicate some "symmetric" properties of OGERs.

**Lemma 2.1** The following are equivalent:

- $an (r, t; s_1, s_2)$ -OGER,
- $an (r, t; s_2, s_1)$ -OGER,
- $a(t,r;s_1,s_2)$ -OGER, and
- $a(t, r; s_2, s_1)$ -OGER.

**Lemma 2.2** There exists an  $(r, t; s_1, s_2)$ -OGER if and only if there exists an  $(s_1, s_2; r, t)$ -OGER.

*Proof.* Suppose  $A = (a_{i,j})$  and  $B = (b_{i,j})$ , where  $1 \le i \le r, 1 \le j \le t$ , form an  $(r, t; s_1, s_2)$ -OGER. Construct two  $s_1 \times s_2$  rectangles  $A' = (a'_{m,n})$  and  $B' = (b'_{m,n})$ , where  $a'_{m,n} = i$  and  $b'_{m,n} = j$  if and only if  $(a_{i,j}, b_{i,j}) = (m, n)$ . It is readily verified that A' and B' form an  $(s_1, s_2; r, t)$ -OGER.

We will make essential use of the Kronecker product. Let  $C = (c_{i,j})$  be an  $r_1 \times t_1$  array, and let  $D = (d_{i,j})$  be an  $r_2 \times t_2$  array. Define an  $r_1r_2 \times t_1t_2$  array  $E = C \bigotimes D = (e_{i,j})$ , where

$$e_{i,j} = (c_{n,q}, d_{m,p}), \text{ for } i = nr_2 + m, j = qt_2 + p, 0 \le m < r_2, 0 \le p < t_2.$$

E is the Kronecker product of C and D.

We now present the three recursive constructions we use.

**Construction 2.3** If there exists a (c, b; c, b)-OGER and a (d, a; a, d)-OGER, then there exists a (cd, ab; ac, bd)-OGER.

*Proof.* We begin with two OGERs. The first is a (c, b; c, b)-OGER consisting of rectangles C and D and the second is a (d, a; a, d)-OGER consisting of rectangles E and F. Now let  $A = C \bigotimes E$  and  $B = D \bigotimes F$ . We prove that A and B are the desired (cd, ab; ac, bd)-OGER.

For the *i*th row of A, where i = nd + m, the elements are  $(c_{n,q}, e_{m,p}), 0 \le q < b, 0 \le p < a$ . Since each symbol in C appears  $\lceil \frac{b}{c} \rceil$  or  $\lfloor \frac{b}{c} \rfloor$  times in a row and each symbol appears exactly once in a row of E, each pair of the symbols appears  $\lceil \frac{b}{c} \rceil$  or  $\lfloor \frac{b}{c} \rfloor$  times in a row of A. In a similar way we can check that conditions 2 and 3 of the definition are satisfied. Finally, it is straightforward to prove that A and B are orthogonal.

**Construction 2.4** If there exists an (m, n; n, m)-OGER, where  $(m, n) \neq (1, 1)$ , then there exists an (2m, 3n; 2n, 3m)-OGER.

*Proof.* First, suppose that  $n \ge 2$ . Suppose  $A = (a_{i,j})$  and  $B = (b_{i,j})$ , where  $1 \le i \le m, 1 \le j \le n$ , are an (m, n; n, m)-OGER. Let  $A_1, A_2$  be two copies of A using two different symbol sets and let  $B_1, B_2, B_3$  be three copies of B using three different symbol sets. For an  $m \times n$  matrix  $X = (x_{i,j})$ , let  $X^1 = (x_{i,j})$ , where  $1 \le i \le m, 1 \le j \le \lceil \frac{n}{3} \rceil, X^2 = (x_{i,j})$ , where  $1 \le i \le m, \lceil \frac{n}{3} \rceil + 1 \le j \le 2\lceil \frac{n}{3} \rceil$ , and the remainder of X as  $X^3$ . Observe that  $X^1$  and  $X^2$  always have the same width.  $X^3$  has the

same width as  $X^1$  and  $X^2$  when  $n \equiv 0 \mod 3$ ; when  $n \not\equiv 0 \mod 3$ ,  $X^3$  is narrower than both  $X^1$  and  $X^2$ .

Construct two  $2m \times 3n$  matrices C and D as follows:

$$C = \begin{bmatrix} A_1 & A_2 & A_1 \\ A_2 & A_1^2, A_1^3, A_1^1 & A_2 \end{bmatrix} \qquad D = \begin{bmatrix} B_1^1, B_2^2, B_3^3 & B_2^1, B_3^2, B_1^3 & B_3^1, B_1^2, B_2^3 \\ B_3^1, B_1^2, B_2^3 & B_3^2, B_1^3, B_1^1 & B_1^1, B_2^2, B_3^3 \end{bmatrix}$$

In the above diagram, commas indicate matrices that are placed side by side.

It is easy to see that C and D form an (2m, 3n; 2n, 3m)-OGER. The only tricky part is to check the alignment of the following subarrays of D (these subarrays will not be perfectly aligned when  $n \neq 0 \mod 3$ ):

$$\begin{array}{c} B_3^2, B_1^3 \\ B_1^3, B_2^1 \end{array}$$

The important point is that there is no overlap of the two occurrences of  $B_1^3$ .

When n = 1, the construction given above does not work. But this does not cause any difficulties. Note that the hypotheses require that m > 1 when n = 1. Using the fact that an (m, 1; 1, m)-OGER is equivalent to a (1, m; m, 1)-OGER (Lemma 2.1), we can construct a (2, 3m; 2m, 3)-OGER by the method described above. By Lemma 2.2, this is equivalent to a (2m, 3; 2, 3m)-OGER.

Similarly, we have the following construction.

**Construction 2.5** If there exists an (m, n; n, m)-OGER, where  $(m, n) \neq (1, 1)$ , then there exists an (3m, 4n; 3n, 4m)-OGER.

*Proof.* Suppose  $A = (a_{i,j})$  and  $B = (b_{i,j})$ , where  $1 \le i \le m, 1 \le j \le n$ , are an (m, n; n, m)-OGER. Let  $A_1, A_2, A_3$  be three copies of A using three different symbol sets and let  $B_1, B_2, B_3, B_4$  be four copies of B using four different symbol sets. For an  $m \times n$  matrix  $X = (x_{i,j})$ , denote  $X = X^1 X^2 X^3$  as in the proof of Construction 2.4.

Construct two  $3m \times 4n$  rectangles C and D as follows:

	$A_1$	$A_2$	$A_3$	$A_1$		$B_1^1, B_2^2, B_3^3$	$B_2^1, B_3^2, B_4^3$	$B_3^1, B_4^2, B_1^3$	$B_4^1, B_1^2, B_2^3$
C =	$A_2^1, A_2^3, A_2^2$	$A_3$	$A_1$	$A_2$	D =	$B_3^1, B_1^3, B_4^2$	$B_4^1, B_1^2, B_2^3$	$B_2^1, B_3^2, B_4^3$	$B_1^1, B_2^2, B_3^3$
	$A_3^2, A_3^3, A_3^1$	$A_1$	$A_2$	$A_3$		$B_2^2, B_3^3, B_1^1$	$B_3^1, B_4^2, B_1^3$	$B_4^1, B_1^2, B_2^3$	$B_2^1, B_3^2, B_4^3$

It is simple to show that C and D form an (3m, 4n; 3n, 4m)-OGER. As in the proof of Lemma 2.4, there are certain subarrays of D that are not perfectly aligned when  $n \neq 0 \mod 3$ :

$B_2^2, B_3^3$
$B_1^3, B_4^2$
$B_3^3, B_1^1$

It is easy to check that there is no overlap of the two occurrences of  $B_3^3$ , nor is there an overlap of  $B_1^3$  and  $B_1^1$ .

The case n = 1 is handled as in Construction 2.4.

**Example 2.6** We illustrate the application of Construction 2.5 with m = 1, n = 4. The following arrays A and B form a (1,4;4,1)-OGER:

$$A = \begin{bmatrix} 1 & 2 & 3 & 4 \end{bmatrix} \qquad \qquad B = \begin{bmatrix} 1 & 1 & 1 & 1 \end{bmatrix}$$

Then  $A^1, A^2, B^1$  and  $B^2$  have width 2, while  $A^3$  and  $B^3$  are empty. We construct C and D, which form a (3, 16; 12, 4)-OGER;

ĺ	1	2	3	4	5	6	7	8	9	a	b	c	1	2	3	4
C =	5	6	7	8	9	a	b	c	1	2	3	4	5	6	7	8
	b	c	9	a	1	2	3	4	5	6	7	8	9	a	b	c
•																
	1	1	2	2	2	2	3	3	3	3	4	4	4	4	1	1
D =	$\frac{1}{3}$	$\frac{1}{3}$	2 $4$	2 4	$\frac{2}{4}$	2 4	3 1	$\frac{3}{1}$	$\frac{3}{2}$	$\frac{3}{2}$	$\frac{4}{3}$	$\frac{4}{3}$	4	4	$\frac{1}{2}$	$\frac{1}{2}$
D =	1 3 2	1 3 2	2 4 1	2 4 1		2 4 3	3 1 4	3 1 4					4 1 2	4 1 2	1 2 3	1 2 3

**Construction 2.7** There exist a (6, 6; 4, 9)-OGER, a (6, 12; 8, 9)-OGER and a (12, 12; 9, 16)-OGER.

*Proof.* These three OGERs are each constructed using a similar technique. For positive integers r and s, define  $c = \operatorname{lcm}(r, s)/r$ . Then define an  $r \times c$  array  $D_{r,s}$  having entries  $d_{i,j} = jr + i \mod s$ ,  $0 \leq j \leq c - 1$ ,  $0 \leq i \leq r - 1$ . Suppose that c|t, and define  $E_{r,t,s}$  to consist of t/c copies of  $D_{r,s}$  placed side by side.

Next, suppose that  $\pi \in (\mathbb{Z}_r)^t$  and construct  $\pi(E_{r,t,s})$  from  $E_{r,t,s}$  by rotating column j of  $E_{r,t,s}$  upwards cyclically by  $\pi(j)$  positions, for  $j = 0, \ldots, t-1$ .

It can be verified that the following arrays form the desired OGERs:

- $\pi(E_{6,6,4})$  and  $\pi(E_{6,6,9})$ , where  $\pi = (0, 0, 1, 1, 2, 2)$ .
- $\pi(E_{6,12,8})$  and  $\pi(E_{6,12,9})$ , where  $\pi = (0, 0, 0, 0, 1, 1, 1, 1, 2, 2, 2, 2)$ .
- $\pi(E_{12,12,9})$  and  $\pi(E_{12,12,16})$ , where  $\pi = (0, 0, 0, 1, 1, 1, 2, 2, 2, 3, 3, 3)$ .

**Example 2.8** We illustrate the construction of a (6,6;4,9)-OGER. First, we depict  $E_{6,6,4}$  and  $E_{6,6,9}$ :

0	2	0	2	0	2
1	3	1	3	1	3
2	0	2	0	2	0
3	1	3	1	3	1
0	2	0	2	0	2

0	6	3	0	6	3
1	7	4	1	7	4
2	8	5	2	8	5
3	0	6	3	0	6
4	1	7	4	1	7
5	2	8	5	2	8

It is not hard to verify that these arrays are orthogonal, and each of them is equitable on columns. Now apply the column rotations specified by  $\pi$  to these two arrays:

						_						
0	2	1	3	2	0		0	6	4	1	8	5
1	3	2	0	3	1		1	7	5	2	0	6
2	0	3	1	0	2		2	8	6	3	1	7
3	1	0	2	1	3		3	0	7	4	2	8
0	2	1	3	0	2		4	1	8	5	6	3
1	3	0	2	1	3		5	2	3	0	7	4

It can be verified that the resulting arrays are now orthogonal, equitable on rows and equitable on columns. Therefore we have a (6, 6; 4, 9)-OGER.

At this point, we are in a position to prove our main result.

**Theorem 2.9** Suppose r, t,  $s_1$  and  $s_2$  are positive integers such that  $rt = s_1s_2$ . Then there exists an  $(r, t; s_1, s_2)$ -OGER if and only if  $(r, t; s_1, s_2) \notin \{(2, 2; 2, 2), (2, 3; 2, 3), (3, 4; 3, 4), (6, 6; 6, 6)\}$ .

*Proof.* Let  $b = \text{gcd}(t, s_2)$ , a = t/b,  $d = s_2/b$  and c = r/d. Then gcd(a, d) = 1. It is clear that a, b and d are integers; we prove now that c is also an integer. Since  $rt = s_1s_2$ , we have

$$c = \frac{r}{d} = \frac{rt}{dt} = \frac{s_1 s_2}{dt} = \frac{s_1 bd}{dba} = \frac{s_1}{a}.$$

On the other hand,

$$\frac{s_1d}{a} = \frac{s_1db}{ab} = \frac{s_1s_2}{ab} = \frac{s_1s_2}{t} = r$$

is an integer. From the fact that gcd(a, d) = 1, it follows that  $c = s_1/a$  is an integer.

Therefore we have that  $(r, t; s_1, s_2) = (cd, ab; ac, bd)$ , where a, b, c and d are positive integers. By Construction 2.3, if there exist a (c, b; c, b)-OGER and a (d, a; a, d)-OGER, then there exists an  $(r, t; s_1, s_2)$ -OGER. So we just need to consider the exceptions from Theorem 1.2.

We consider three cases, as follows.

- 1. There is a (c, b; c, b)-OGER, where c and b are not both equal to one, but a (d, a; a, d)-OGER does not exist. For (d, a; a, d) = (2, 2; 2, 2) or (6, 6; 6, 6), the designs are constructed in Theorem 1.2. For (d, a; a, d) = (2, 3; 3, 2) or (3, 4; 4, 3), the designs are constructed in Constructions 2.4 and 2.5.
- 2. There is a (d, a; a, d)-OGER, where d and a are not both equal to one, but a (c, b; c, b)-OGER does not exist. This is equivalent to case 1, by Lemma 2.1.
- 3. Both (c, b; c, b)-OGER and (d, a; a, d)-OGER do not exist. When one of the missing OGERs is of type (2, 2; 2, 2) or (6, 6; 6, 6), then the designs are constructed in Theorem 1.2. So we just need to consider the exceptions (2, 3; 2, 3) and (3, 4; 3, 4). Using Lemmas 2.2 and 2.1, there are three types of OGERs that we need to construct: (6, 6; 4, 9), (6, 12; 8, 9), and (12, 12; 9, 16). These were handled in Construction 2.7.

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